

# Econ 701: Answers to Problem Set 1

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1. First, let's show that  $\succ$  is asymmetric if and only if  $\succeq$  is complete. By definition of completeness,  $\succeq$  is complete if and only if for every  $x$  and  $y$ , either  $x \succeq y$  or  $y \succeq x$ . Hence by definition of  $\succeq$ ,  $\succeq$  is complete iff for every  $x$  and  $y$ , either  $x \not\succeq y$  or  $y \not\succeq x$ . But this is exactly the definition of  $\succ$  being asymmetric.

Next, let's show that  $\succ$  is negatively transitive if and only if  $\succeq$  is transitive. By definition,  $\succ$  is negatively transitive iff whenever  $x \succ y$ , we have either  $z \succ y$  or  $x \succ z$ . Taking the contrapositive, this holds iff whenever  $z \not\succeq y$  and  $x \not\succeq z$ , we have  $x \not\succeq y$ . By definition of  $\succeq$ , this is true iff whenever  $y \succeq z$  and  $z \succeq x$ , we have  $y \succeq x$ . This is exactly the definition of  $\succeq$  being transitive.

2. Assume  $\succeq$  is transitive. Suppose  $x \succeq y$  and  $y \succ z$ . By definition, the second statement says  $y \succeq z$  and  $z \not\succeq y$ . By transitivity of  $\succeq$ ,  $x \succeq y$  and  $y \succeq z$  implies  $x \succeq z$ . So we need to show that  $z \not\succeq x$ . Suppose not. Then  $z \succeq x$ . Then we have  $z \succeq x$  and  $x \succeq y$ , so transitivity of  $\succeq$  implies  $z \succeq y$ . But we said above that  $z \not\succeq y$ , so this is impossible. Hence  $z \not\succeq x$  so  $x \succ z$ .

3. This procedure can violate WARP but doesn't necessarily. For example, if  $\delta$  is smaller than the smallest possible utility difference, then the procedure won't violate WARP. In other words, if whenever  $u(x) > u(y)$ , we actually have  $u(x) > u(y) + \delta$ , then the  $\delta$  is irrelevant and this agent is really just maximizing  $u$ . So he will necessarily satisfy WARP.

For an example violating WARP, suppose  $u(x_1) = 0$ ,  $u(x_2) = 9$ ,  $u(x_3) = 10$ , and  $\delta = 1.1$ . If  $B = \{x_1, x_2, x_3\}$ , then  $x^* = x_1$ ,  $u(x^*) = 0$ , and  $\bar{u} = 10$ . Since  $\bar{u} > u(x^*) + \delta$ , the decision maker chooses  $x_3$ . Now suppose  $\bar{B} = \{x_2, x_3\}$ . WARP would require  $x_3$  to still be chosen. However, now we have  $x^* = x_2$ , so  $u(x^*) = 9$ . Hence we no longer have  $\bar{u} > u(x^*) + \delta$ , so the decision maker chooses  $x_2$ .

4. This one must violate WARP. Simply note that  $C(\{x_1, x_3\}) = \{x_1\}$  and  $C(\{x_1, x_2, x_3\}) = \{x_3\}$ . Again,  $x_1$  and  $x_3$  are in both sets, only  $x_1$  is chosen from the first, and only  $x_3$  is chosen from the second, a violation of WARP.

5. One point I should have clarified: When I stated the alternative definition of WARP

in class, I stated it as an “iff.” It’s really only an “iff” in the case where  $\mathcal{B}$  is rich. To illustrate this, I’ll write out a longer answer than you needed to give which explains this point.

So suppose our first definition of WARP is satisfied. That is, if  $x, y \in B \cap \bar{B}$ ,  $x \in C(B)$ , and  $y \in C(\bar{B})$ , then we have  $x \in C(\bar{B})$ . Now let’s show the second definition (using “iff”) must hold. So I want to show that  $x \not\prec_C^* y$  iff  $y$  is strictly revealed preferred to  $x$ . So we need to rule out the possibility that the first holds but not the second or that the second holds but not the first. Let’s take these one at a time. Suppose  $x \not\prec_C^* y$  but  $y$  is not strictly revealed preferred to  $x$ . Here we need richness.  $y$  not strictly revealed preferred to  $x$  says that  $\{y\} \neq C(\{x, y\})$ . Note that this uses richness since  $\{x, y\}$  might not be in  $\mathcal{B}$  otherwise. Since  $C(\{x, y\})$  is nonempty, this tells us  $x \in C(\{x, y\})$ , so  $x \succeq_C^* y$ , a contradiction. So now rule out the second possibility. Suppose  $x \succeq_C^* y$  but  $y$  is strictly revealed preferred to  $x$ . By definition, the first statement says that there is some set  $B$  such that  $x, y \in B$  and  $x \in C(B)$ . The second statement says, by definition, that there is some set  $\bar{B}$  with  $x, y \in \bar{B}$ ,  $y \in C(\bar{B})$ , and  $x \notin C(\bar{B})$ . But then  $x, y \in B \cap \bar{B}$ ,  $x \in C(B)$ , and  $y \in C(\bar{B})$ , so we must have  $x \in C(\bar{B})$ , a contradiction. Note that this part does *not* use richness, only the definitions.

In short, our first definition of WARP implies the second with the “iff” if richness holds and implies the second phrased as “if  $x \succeq_C^* y$ , then  $y$  is not strictly revealed preferred to  $x$ .”

For the converse, suppose WARP holds according to the *weaker* version of the second definition — that is, the one that doesn’t have an “iff.” Suppose  $x \succeq_C^* y$ . By the second definition, this means that  $y$  is not strictly revealed preferred to  $x$ . That is, we have that there is some  $B$  with  $x, y \in B$  and  $x \in C(B)$  and that for any  $\bar{B}$  such that  $x, y \in \bar{B}$  and  $y \in C(\bar{B})$ , we also have  $x \in C(\bar{B})$ . But this is our first definition. Note again that this does not use richness.

6. I will show that the function

$$u(x) = \#\{y \in X \mid x \succ y\}$$

represents  $\succeq$  where  $\#$  denotes the cardinality of a set (in the finite case, the number of items in the set). So suppose that  $u(x) \geq u(x')$ . I wish to show that this implies that  $x \succeq x'$ . Suppose not. Then, by definition,  $x' \succ x$ . For simplicity, let

$$A = \{y \in X \mid x \succ y\}$$

and

$$B = \{y \in X \mid x' \succ y\}.$$

Consider any  $z \in A$ . By the definition of  $A$ ,  $x \succ z$ . By hypothesis,  $x' \succ x$  so, by transitivity,  $x' \succ z$ . This implies that  $z \in B$ . By asymmetry, we cannot have  $x \succ x$ , so

we must have  $x \notin A$ . Thus  $B$  contains every element of  $A$  plus at least one thing —  $x$  — which is not in  $A$ . But then the definition of the function  $u(x)$  implies that we must have  $u(x') > u(x)$ , a contradiction. Hence we see that  $u(x) \geq u(x')$  implies  $x \succeq x'$ . To complete the proof, we need to show that  $x \succeq x'$  implies  $u(x) \geq u(x')$ . This is easily shown by a slight variation on the argument given above.

7. Suppose we have such a representation. First, show  $\succ$  is asymmetric. Suppose, to the contrary, that we have  $x$  and  $y$  with  $x \succ y$  and  $y \succ x$ . Then we must have  $u(x) > u(y) + \delta$  and  $u(y) > u(x) + \delta$ . Substituting the second into the right-hand side of the first, we get  $u(x) > u(x) + 2\delta$  or  $\delta < 0$ , a contradiction. Hence  $\succ$  is asymmetric. To show  $\succ$  is transitive, suppose we have  $x \succ y$  and  $y \succ z$ . Then  $u(x) > u(y) + \delta$  and  $u(y) > u(z) + \delta$ . Substituting the second into the right-hand side of the first gives  $u(x) > u(z) + 2\delta$ . Since  $\delta > 0$ , then,  $u(x) > u(z) + 2\delta > u(z) + \delta$ . Hence  $x \succ z$ , so  $\succ$  is transitive.

For the next property, suppose  $w \succ x$  and  $y \succ z$ . Then  $u(w) > u(x) + \delta$  and  $u(y) > u(z) + \delta$ . Suppose we do not have  $w \succ z$ . Then  $u(w) \leq u(z) + \delta$ . Hence  $u(y) > u(z) + \delta \geq u(w) > u(x) + \delta$ , so  $y \succ x$ .

For the final property, suppose  $w \succ x$  and  $x \succ y$ . Then  $u(w) > u(x) + \delta$  and  $u(x) > u(y) + \delta$ . Suppose we do not have  $w \succ y$ . Then  $u(w) \leq u(y) + \delta$ . Hence  $u(w) > u(x) + \delta > u(y) + 2\delta \geq u(y) + \delta$ , so  $w \succ y$ .